Feasibility to replace Interprocess Communication by the Message Passing Interface in microkernel contexts

> René Drolshagen Lasse Löffler

Overview

- Introduction
- Message Passing Interface (MPI)
- Interprocess Communication (IPC)
- Comparison
- Feasibility
- Conclusions

Introduction

- IPC is the communication standard in the microkernel context
- IPC was introduced 35 years ago
- Since this day IPC was only improved and never replaced
- This presentation will do a comparison of both

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History, concept and functionality of the

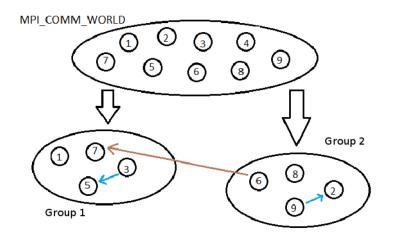
MESSAGE PASSING INTERFACE

- MPI is a specification which tries to solve the problem between portability, efficiency and functionality
- Was developed for distributed shared memory (DSM) architectures
- As trends changed towards NUMA and NoRMA this specification was adapted

- Today MPI could handle all architectures seamlessly and transparently
- Reasons for using MPI:
 - Portability
 - Standardization
 - Performance
 - Functionality

- Communication Methods in MPI:
 - Blocking send / blocking receive
 - Non-blocking send / non-blocking receive
 - Combined send and receive
- Messages are received in the order they were sent by guarantee

- One defined Communicator per Group
 - Handles the whole communication
- Can be compared with the Clans and Chiefs Model



History, concept and functionality of

INTERPROCESS COMMUNICATION

Micro-vs. Monolithic-Kernel

- Operating monolithic kernels are divided in user- and kernelspace
- Issue
 - all the basic services are running in privileged mode
- Idea of the microkernel
 - Moving all the basic services from kernel- to userspace

Idea and usecase of IPC

- Problem
 - Communication between service inside the user space
- Solution
 - Interprocess Communication (IPC) via the microkernels kernelspace

Functionality of IPC

- Passing unbuffered messages between services
- Two fundamental operations
 - send
 - receive
- Both operations can be executed
 - Blocking
 - Non-blocking

Interprocess Communication (IPC)

- IPC needs a high performance
 - Passing thousands of messages
- Also used inside the kernel to handle
 - Interrupts
 - Memory-Managing
- IPC Permissions / Addresses managed via Capabilities

COMPARISON

MPI vs. IPC

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Comparison

- MPI allows n-to-1 and 1-to-n operations
- Asynchron MPI send operations could possibly sleep forever
- IPC has a defined message structure
- MPI messages are defined by the data which is passed inside the message
- MPI has the ability of process groups
- IPC is much faster (!)

Replacing IPC by MPI

FEASIBILITY

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Feasibility

• **Problem:** No Capability feature in MPI

- Every process could communicate with each other
- Solution:
 - Group processes together which are allowed to communicate
 - Only communication inside groups is allowed

• Example:

- Group 1: process A and process B
- Group 2: process A and process C

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Feasibility

• **Problem:** No Capability feature in MPI

- Every process could communicate with each other
- Solution:
 - The communicator of a group has the ability to check whether a communication is allowed
 - Implement a kind of rights table

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Feasibility

- Problem: Denial-of-Service
 - A malicious user-level program could send a unlimited amout of non-blocking messages

• Solution:

• Forbid non-blocking messages

Feasibility

- **Problem:** Starving by blocking messages
 - A performed blocking send operation could starve a process, if the receiver never appears

• Solution:

- a. Implement a timer in messages
- b. Give the communicator the ability to break down the send operation
 - Causes dramativ overhead and implementation work (!)

Feasibility

- **Problem:** Performance and amount of Code
 - MPI has got a huge amount of Code
 - The performance of MPI compared to IPC
- Solution:
 - Not available

THE CONCLUSION

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Conclusion

- Replacing IPC by MPI is not advisable
- Some problems could easily be solved
- The Denial-of-Service attack is a critical issue
- There is a microkernel with MPI, which is called PARAM9000

THANKS FOR YOUR ATTENTION

Questions?

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