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# User-level scheduling

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# What's User-level scheduling

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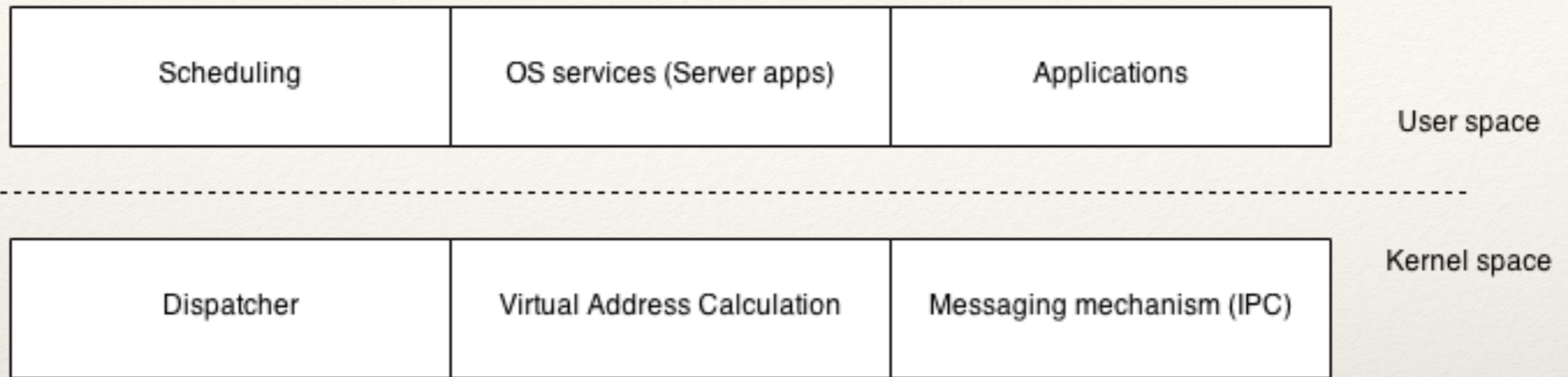
- ❖ Export scheduling policies from the kernel in the User-level
- ❖ Based on the idea of microkernel based operating systems

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# Microkernel based operating systems

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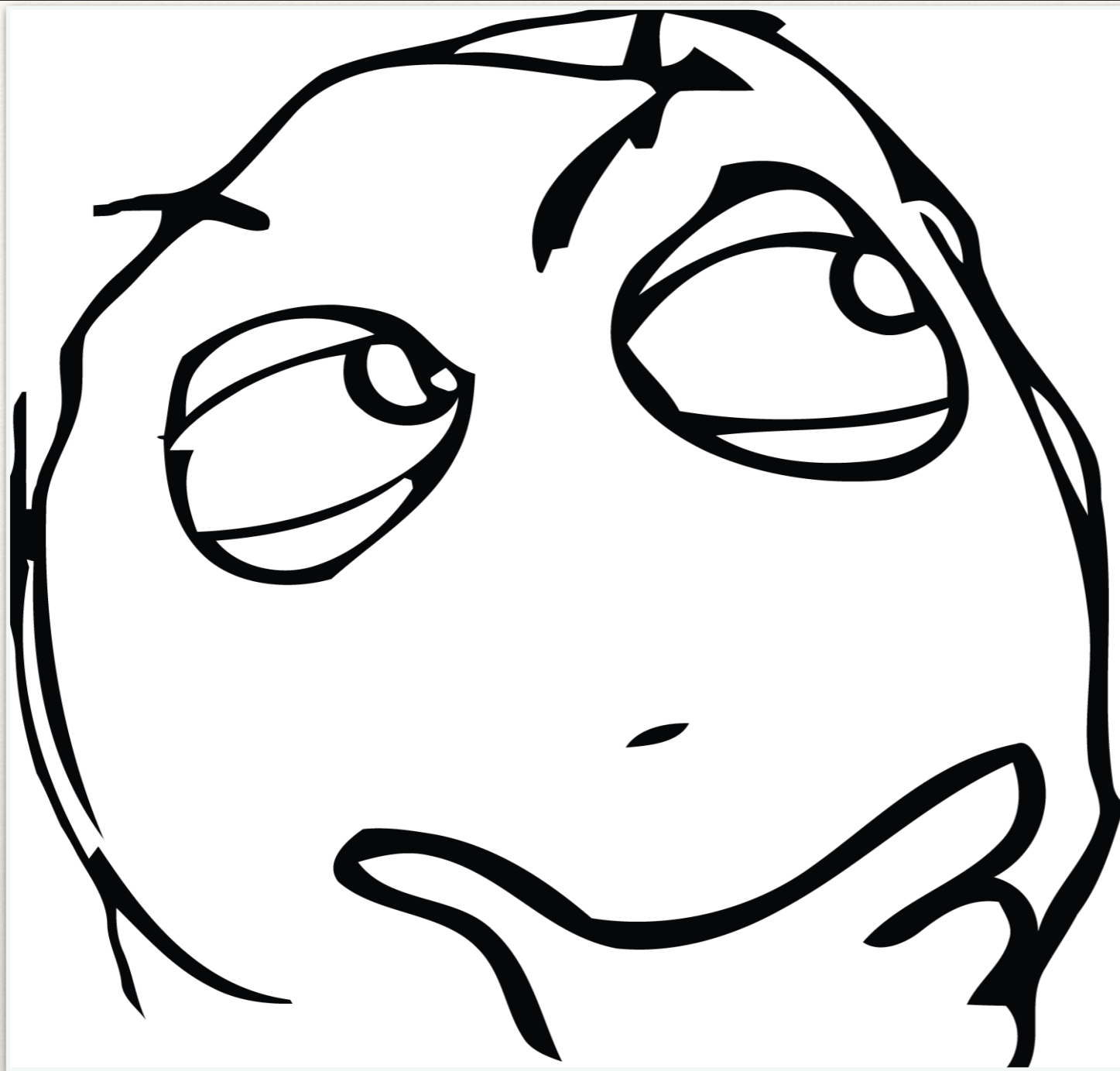
- ❖ Minimize the kernel part of the operating system
- ❖ For more modularity, flexibility and small "Trusted Computing Base"
- ❖ Just include scheduling mechanism, address calculation and a messaging service (IPC)
- ❖ All resource-management policies have to be implemented at user-level as server applications



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Idea of user-level  
scheduling

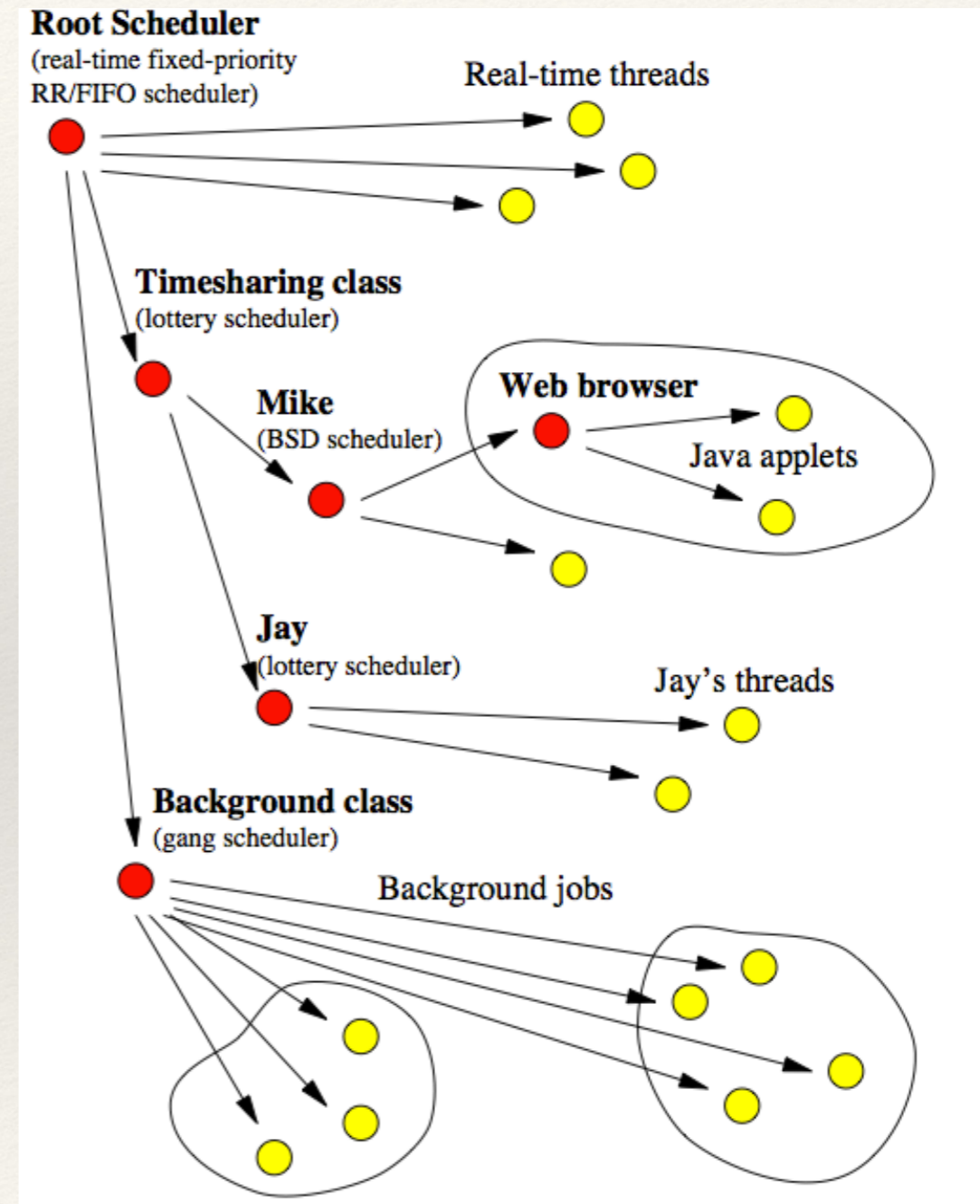
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Why using user-level scheduling?

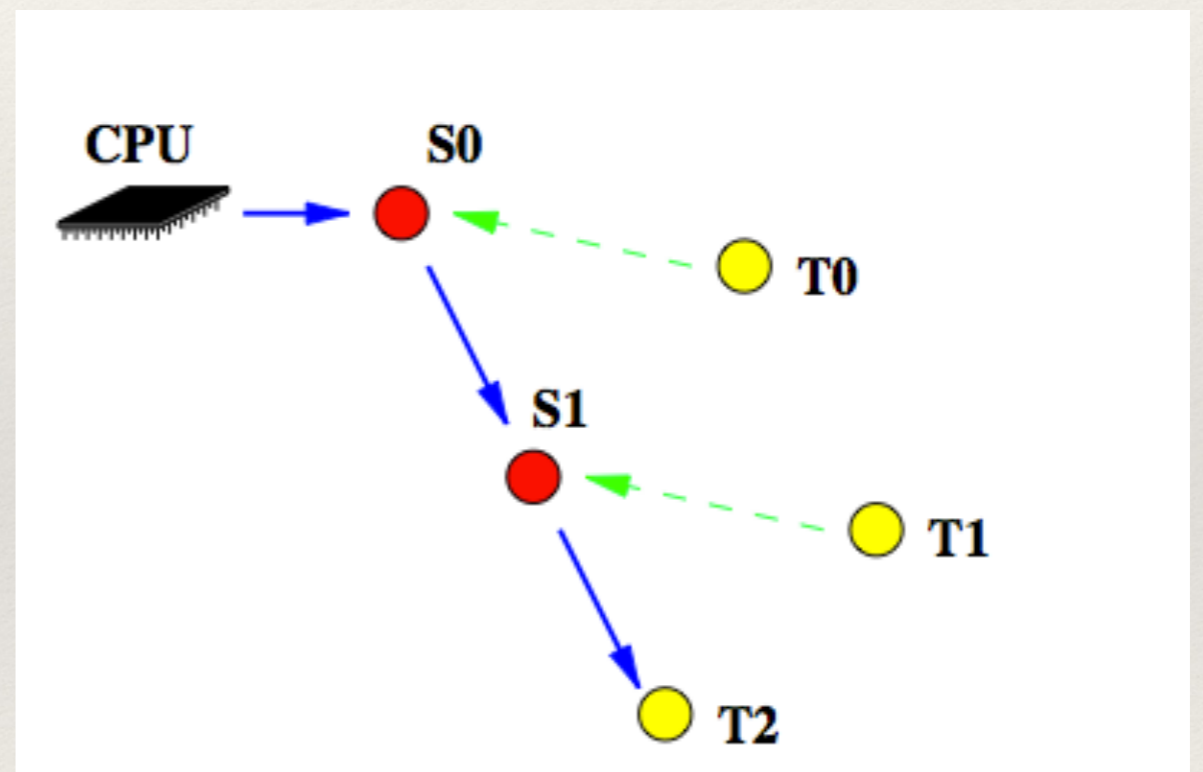
# CPU Inheritance Scheduling

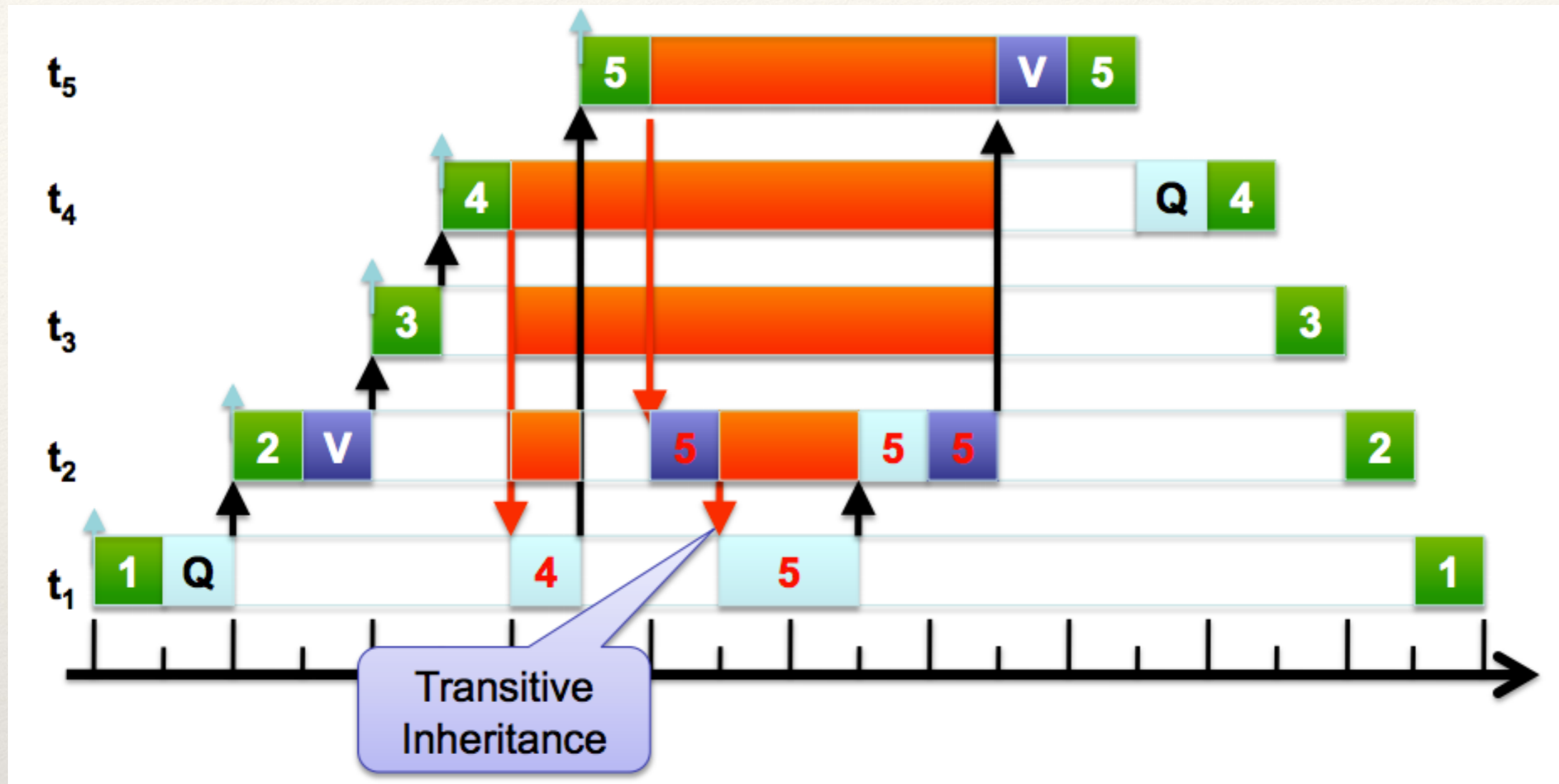
- ❖ Traditionally scheduling is on low level: by kernel scheduler or by user-level thread packages
- ❖ New approach: higher-level threads donate CPU+resources to others
- ❖ "Inheritance": ability to donate and request (virtual) CPU time between threads
- ❖ Client threads can act as scheduler threads for others
- ❖ Root scheduler owns real CPU time



# CPU Usage Accounting

- ❖ Statistical accounting
- ❖ Time stamp-based accounting
- ❖ Directly implement by root schedulers
- ❖ Virtual time information for clients of other schedulers



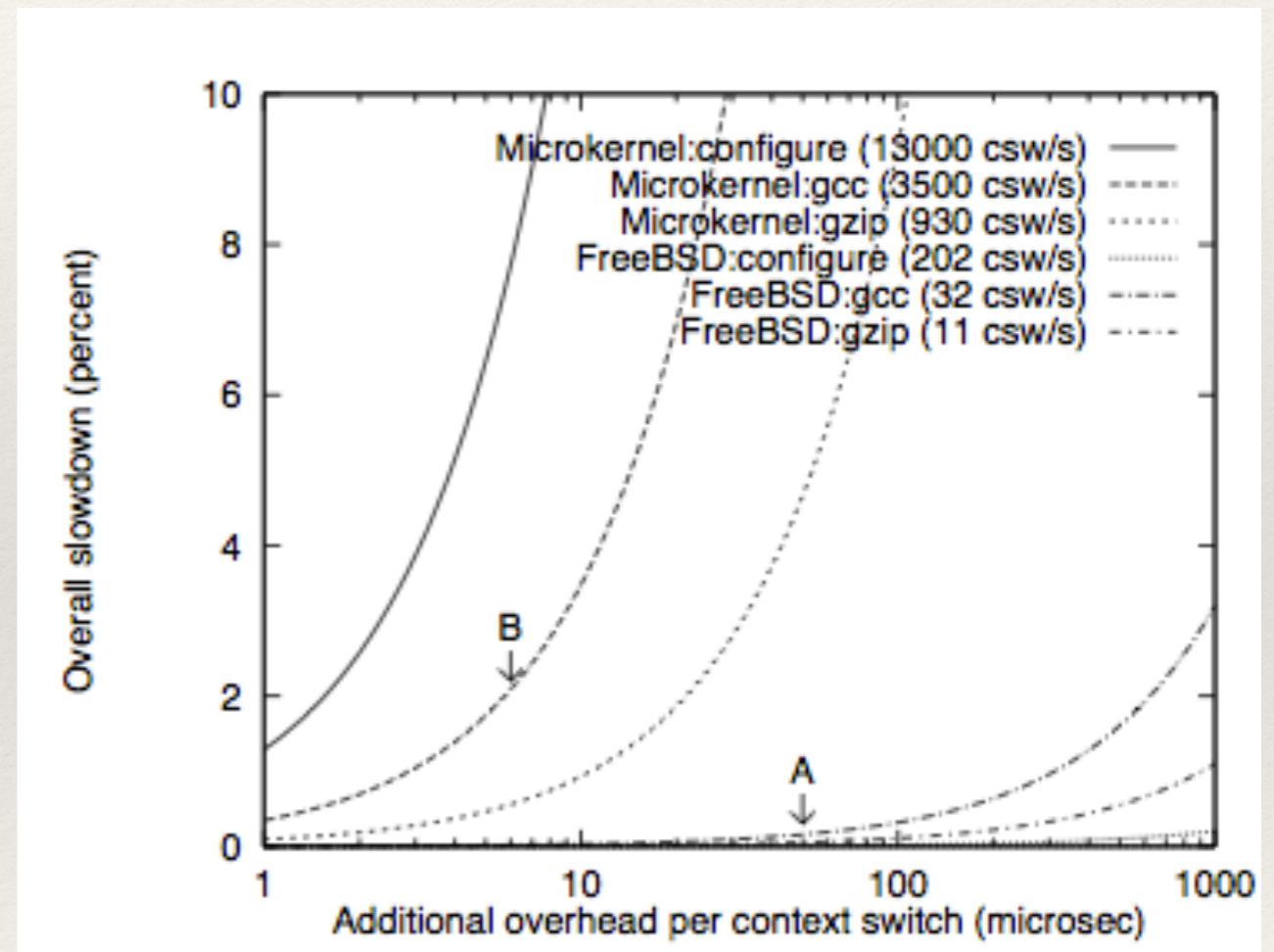


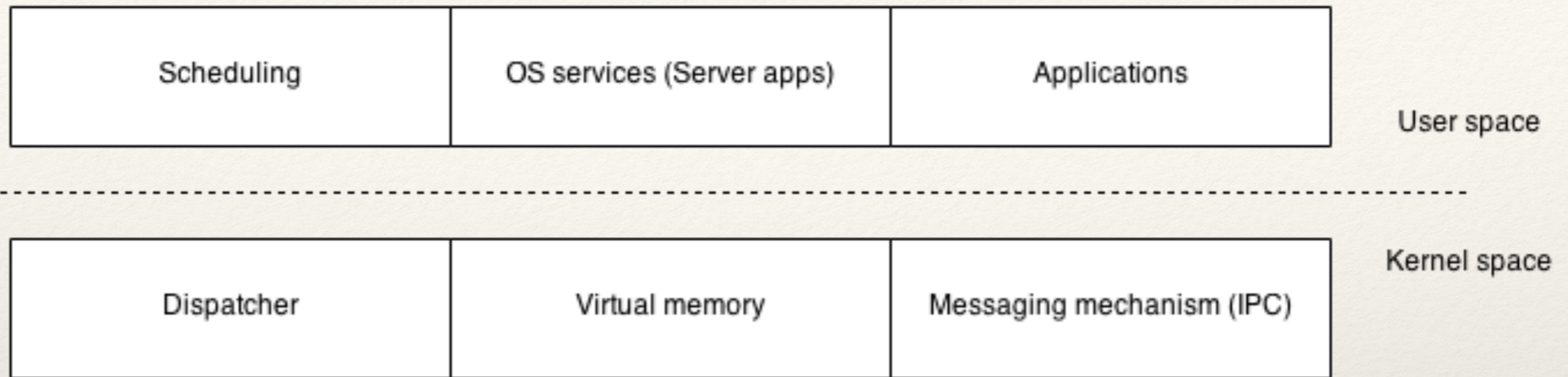
Priority inversion



# Overhead

- ❖ Is it efficient enough in practice?
- ❖ Two additional sources of overhead
  - ❖ Caused by dispatcher
  - ❖ and add. context switches





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*View: Exokernel*

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# Conclusion

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- ❖ CPU inheritance scheduling has low overhead
  - ❖ All threads in one address space
  - ❖ A threads has access to memory of all applications
- ❖ Is it possible to use the concept of CPU inheritance scheduling for more address spaces?

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# That's all Folks

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❖ Thanks for your Attention!