

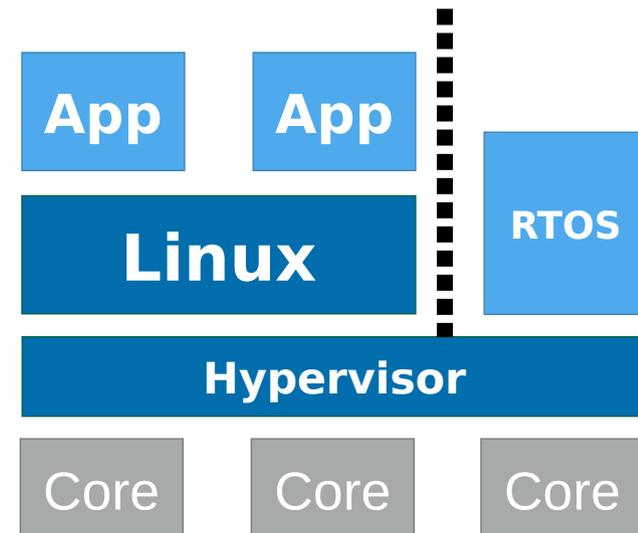
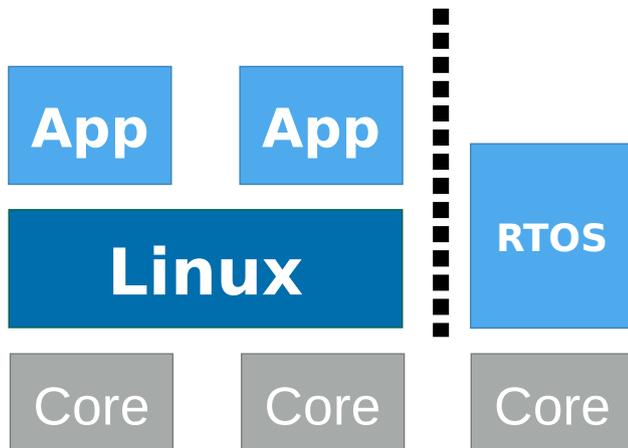
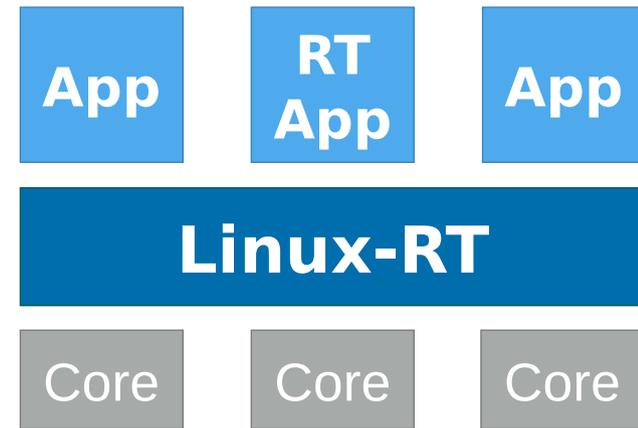
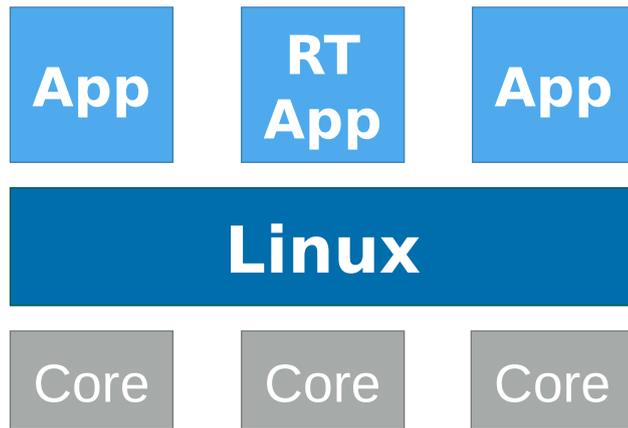
Combining Predictable Execution with Full- Featured Commodity Systems

Adam Lackorzynski, Carsten Weinhold, Hermann Härtig

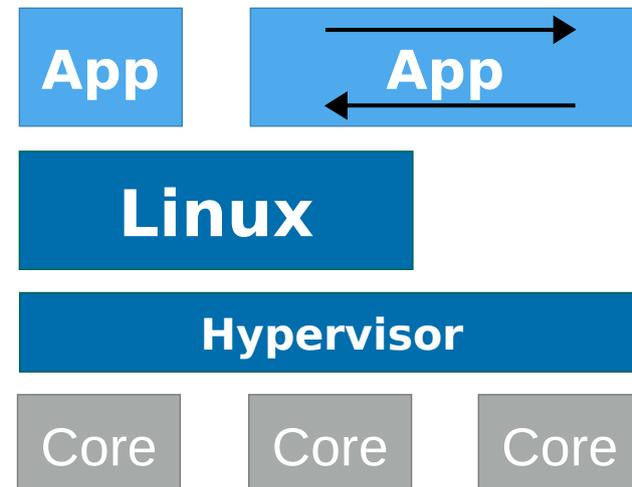
Jul 5th 2016, OSPERT Workshop

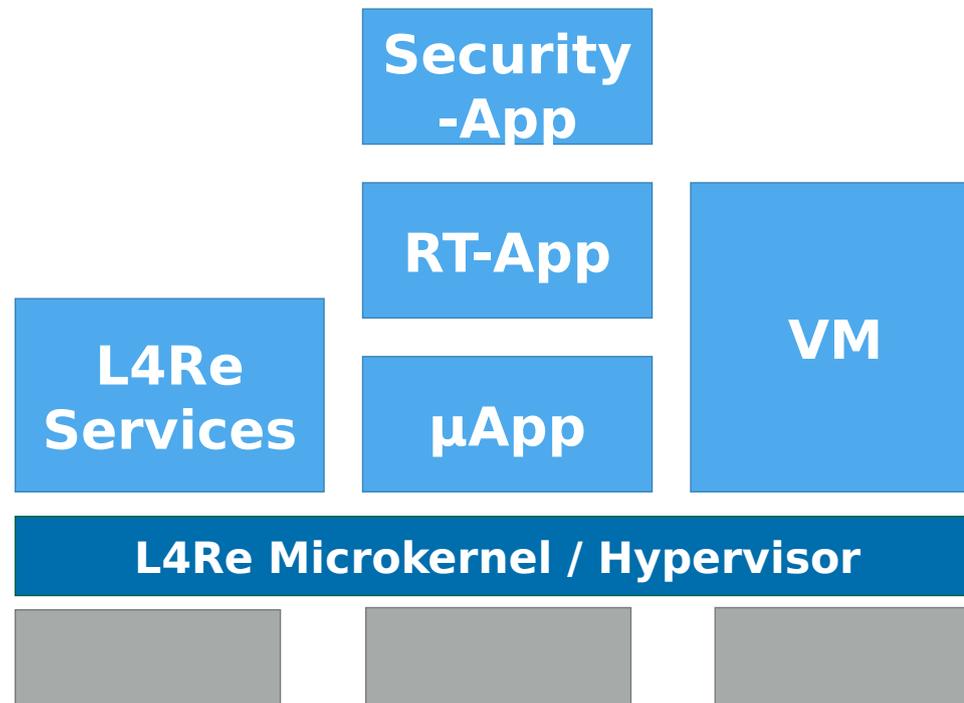


Combining a GPOS and Real-Time

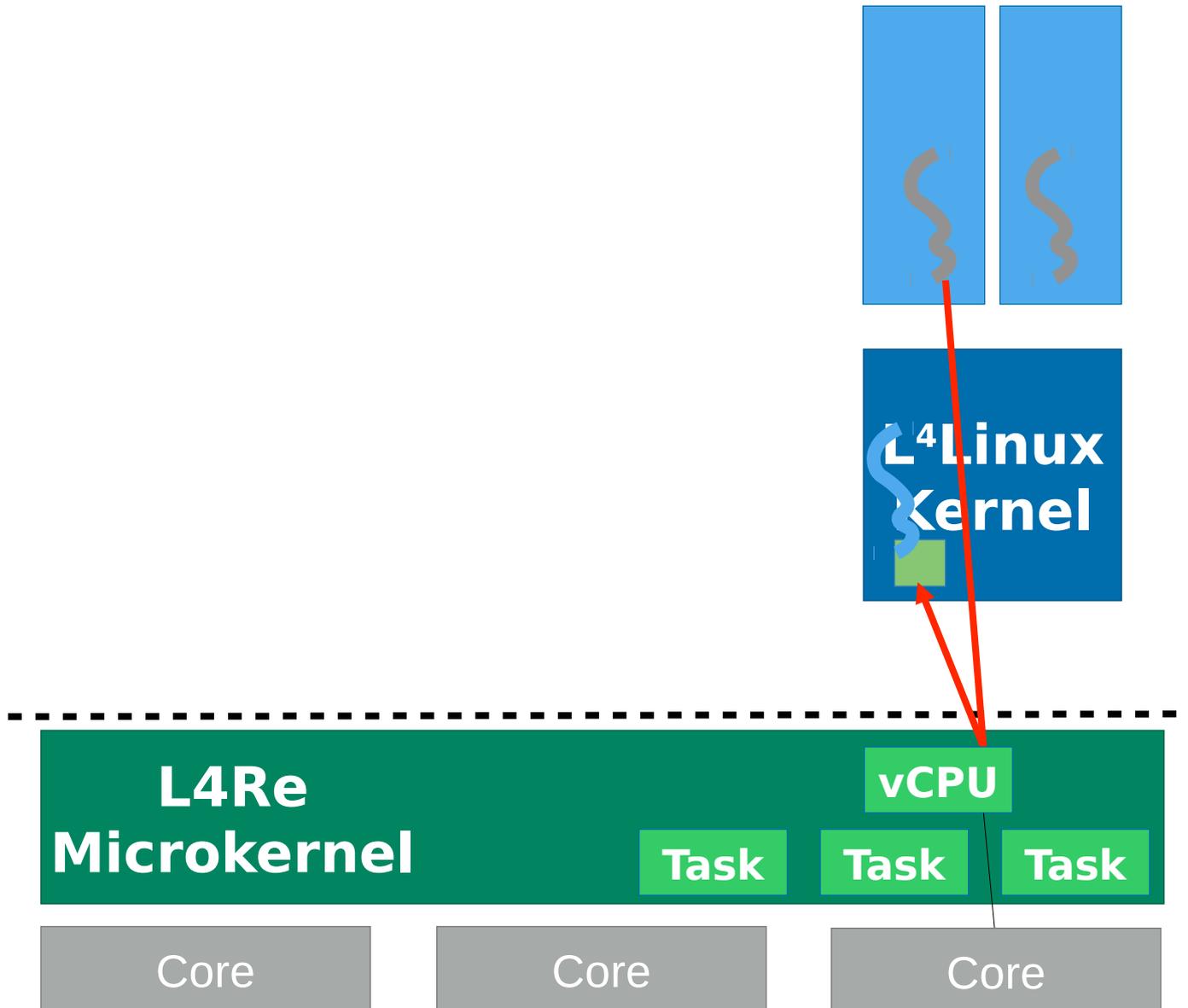


- Use GPOS (Linux) for non-RT parts of an RT application
- Run RT part outside of Linux
 - Separate „noisy“ Linux from sensitive parts
- Best of both worlds:
 - Use Linux for features
 - Run RT outside of Linux
- No gap to overcome
 - No application splitting
 - No proxying

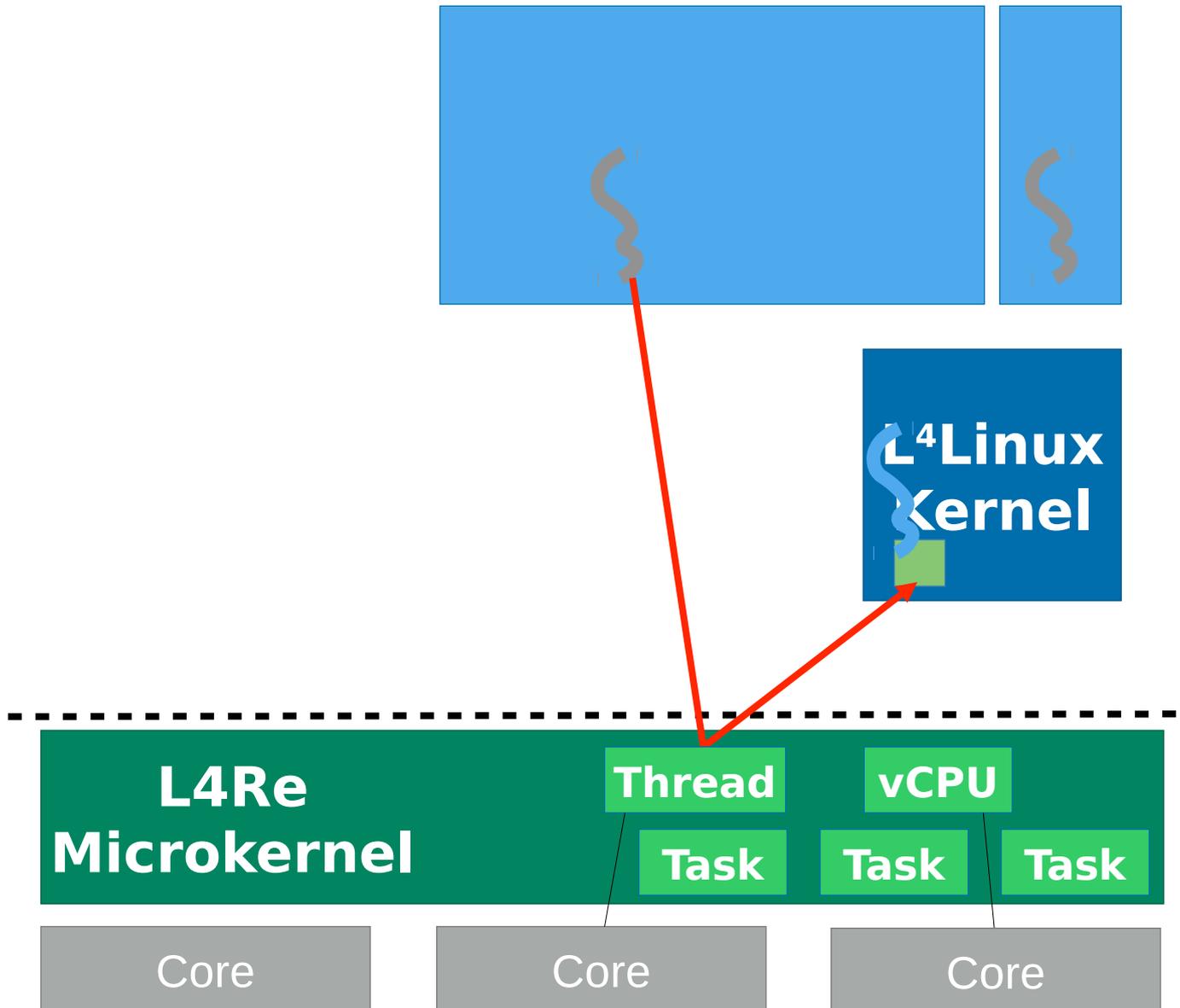




- Hardware-assisted virtualization
 - Intel VT-x, AMDs SVM, ARM VE, MIPS VZ
- Para-Virtualization
 - No hardware support for Virtualization required
 - L⁴Linux, FreeRTOS, L4OpenBSD, ...
 - Tight integration...



Detaching Linux Threads



- Thread of an application can run undisturbed from Linux activities
- Program can execute L4 system calls directly, without involving the Linux kernel
- Multi-threaded applications can run detached and normal in parallel

FWQ (Fixed Work Quantum)

- Measure OS Noise
- Measure execution time jitter

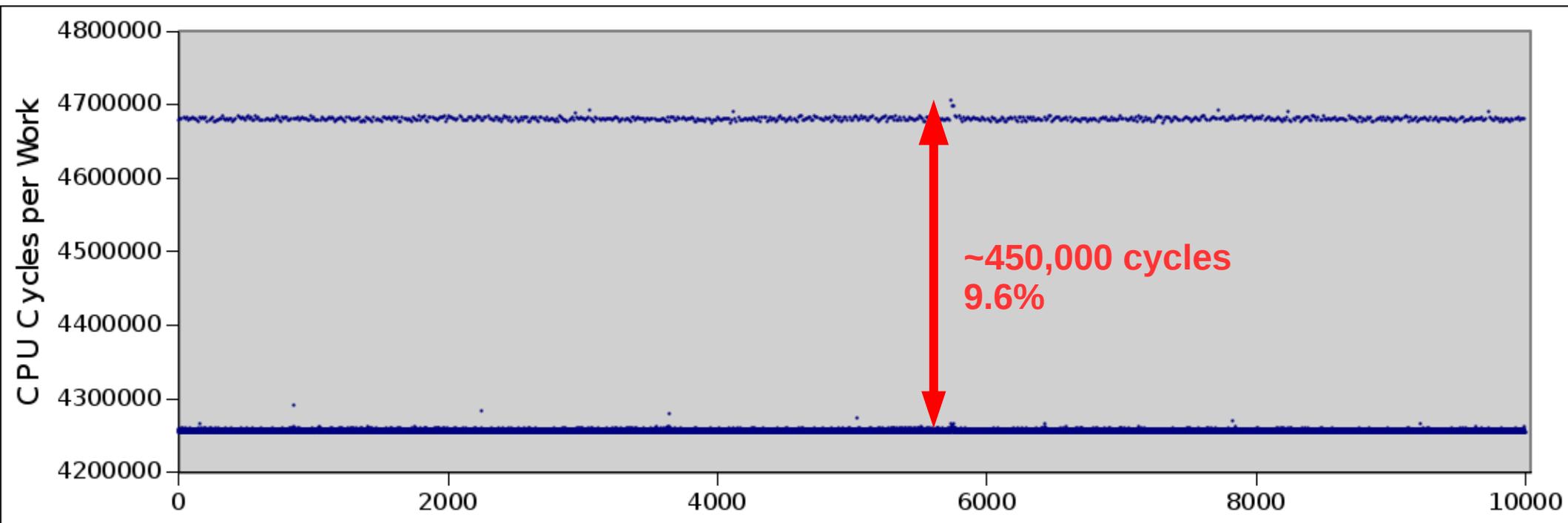
Control loop hybrid application

3 Workloads:

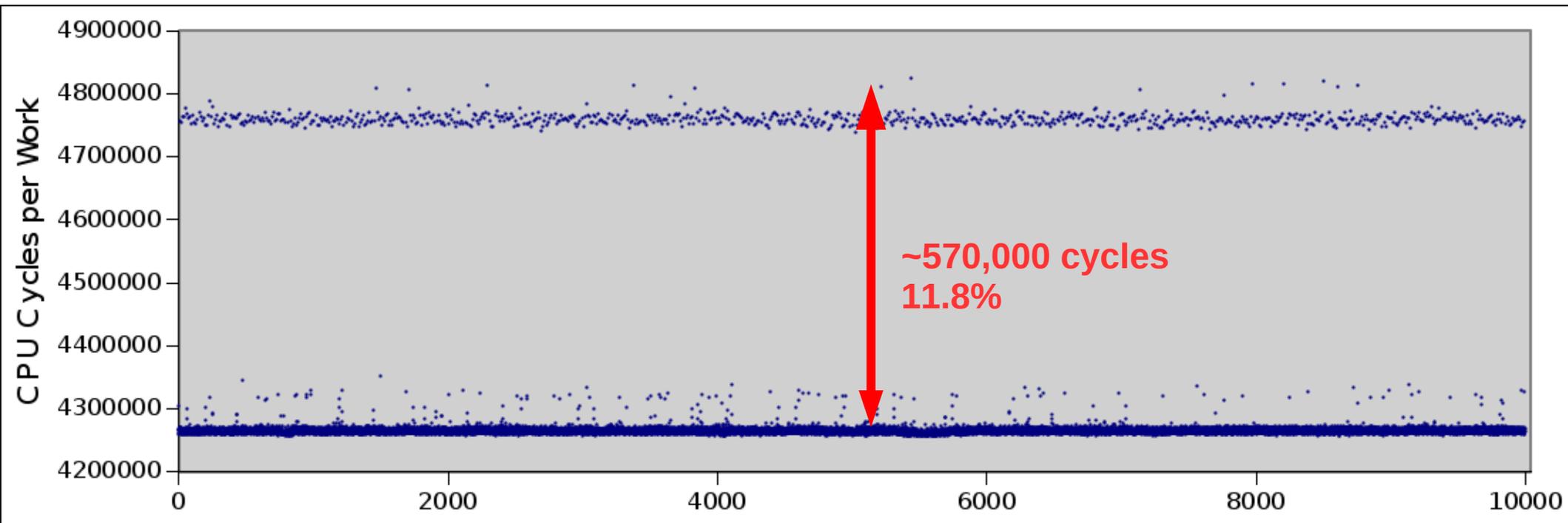
- Idle
- Build: Building a Linux kernel
- IO: Stressing disk and network

Intel Core™ i7-4770, Quad-Core at **2993MHz**

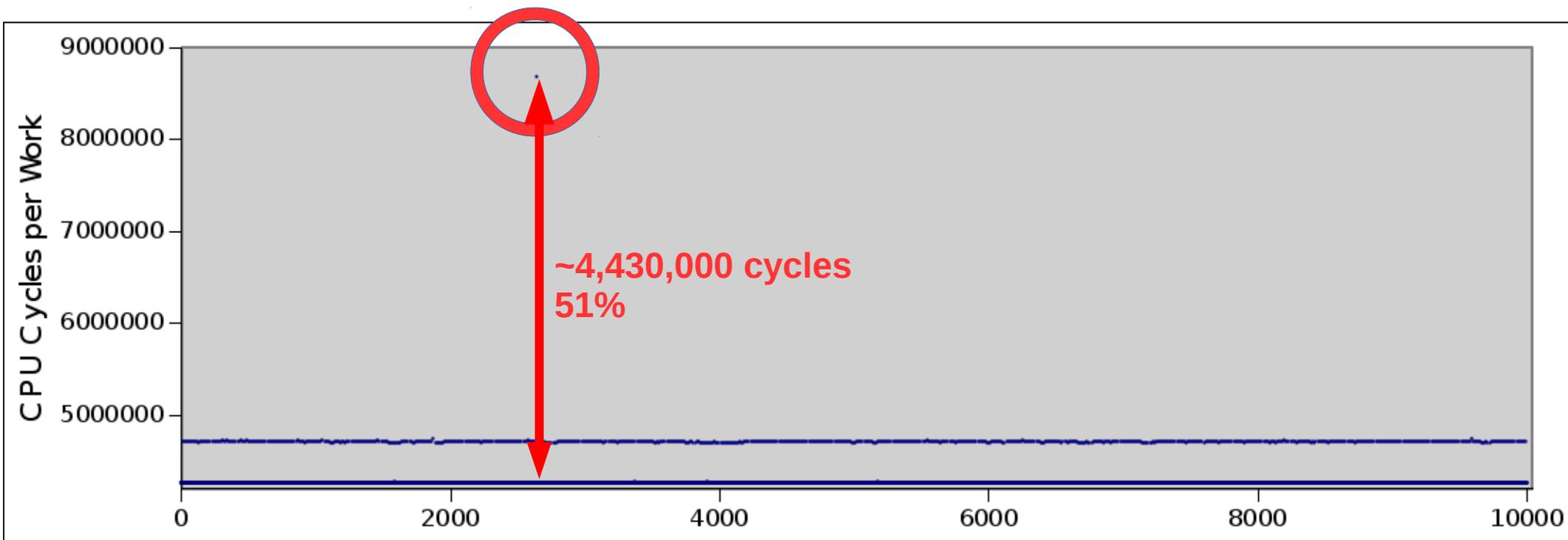
- Linux idle



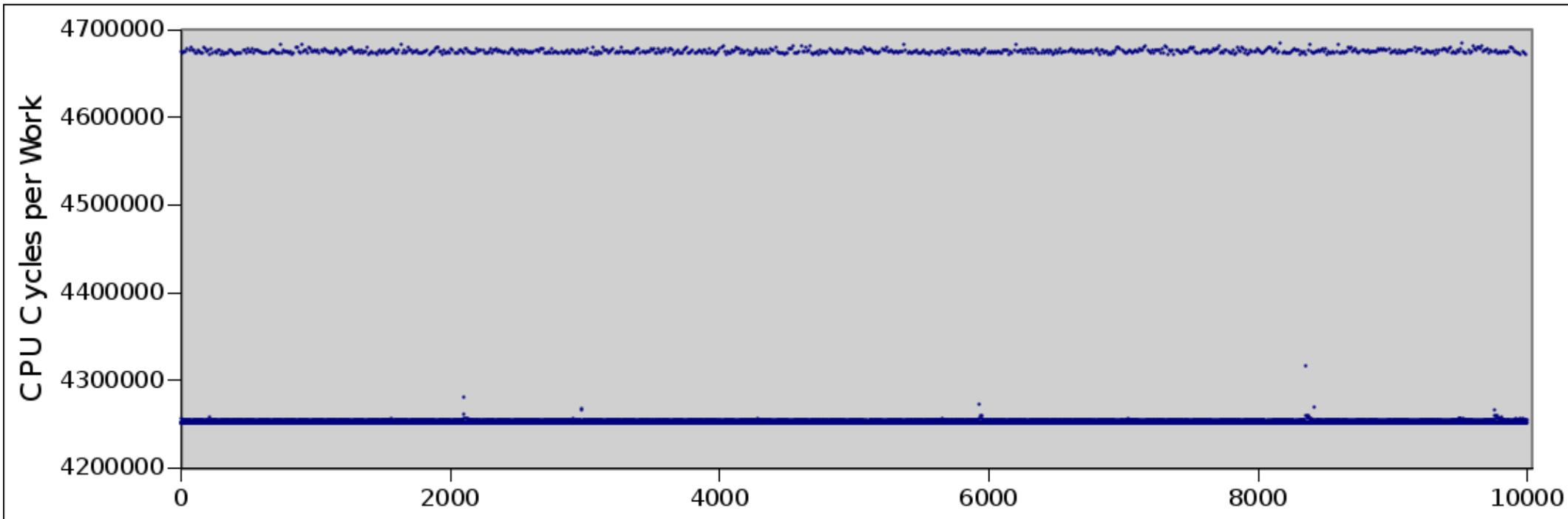
- Linux build load



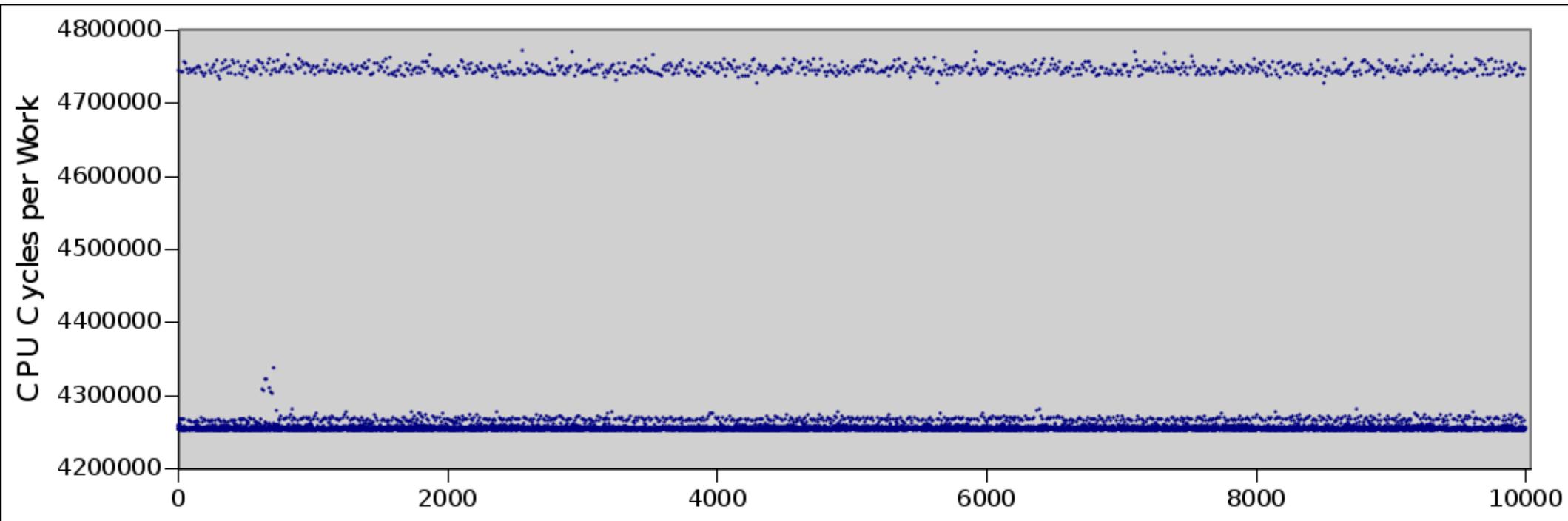
- Linux I/O load



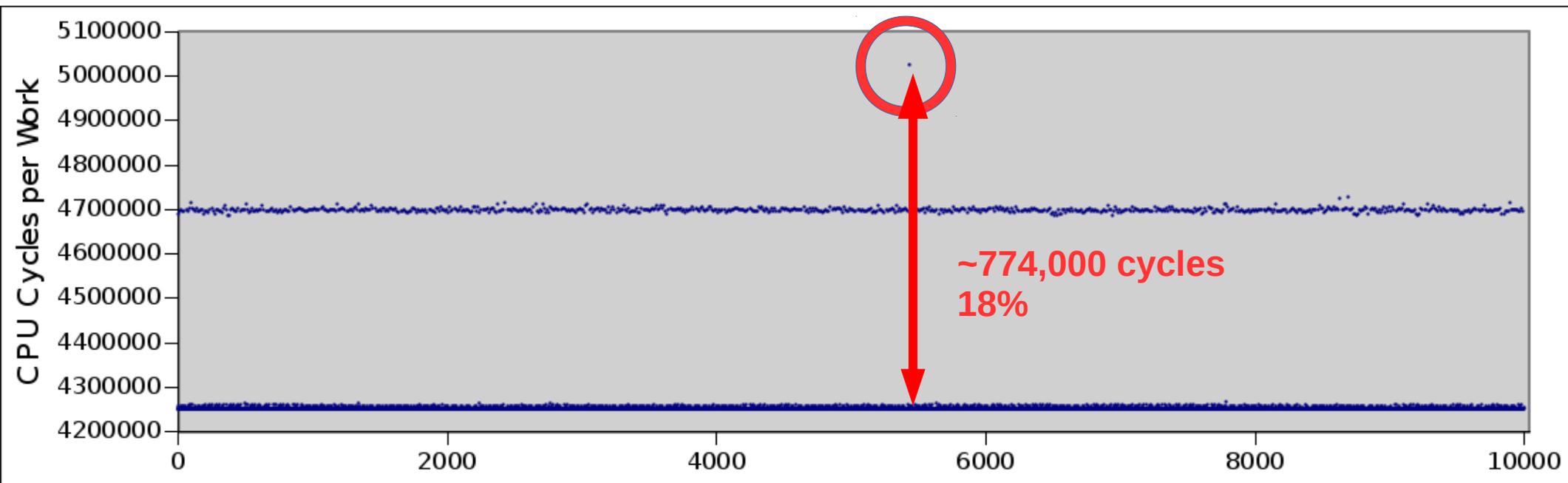
- Linux PREEMPT Idle



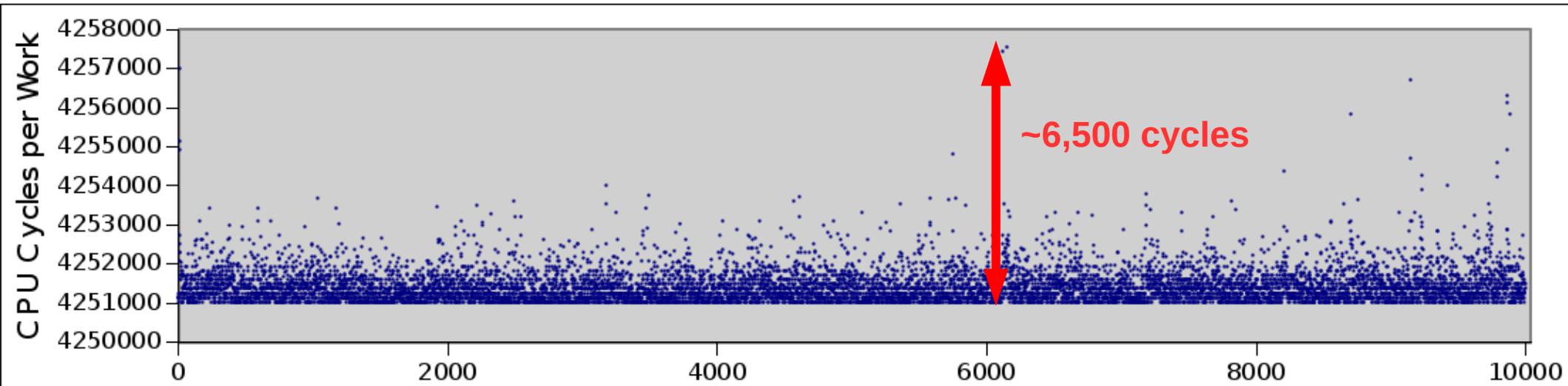
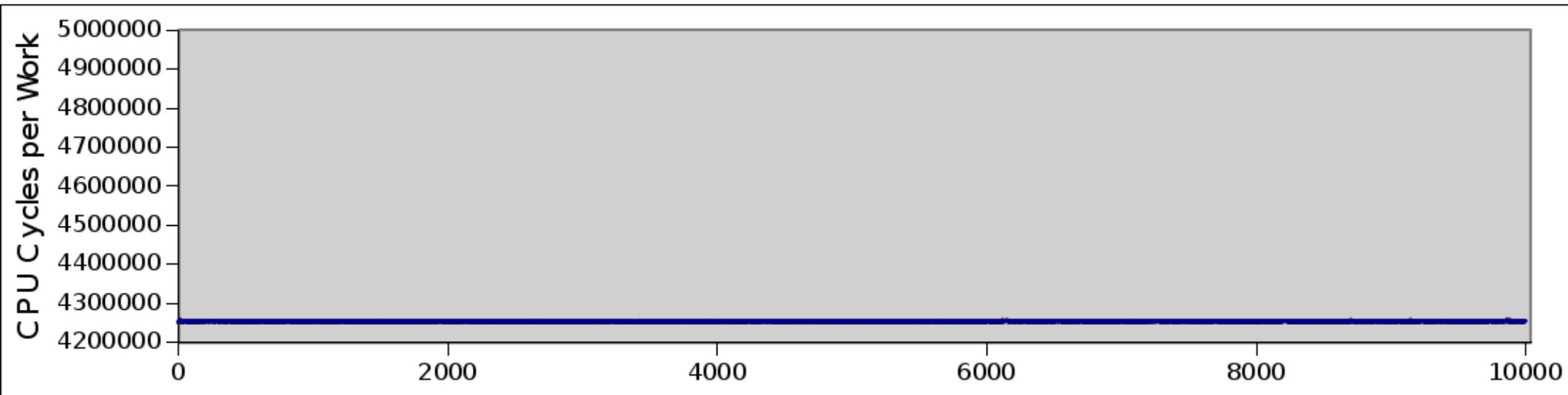
- Linux PREEMPT build load



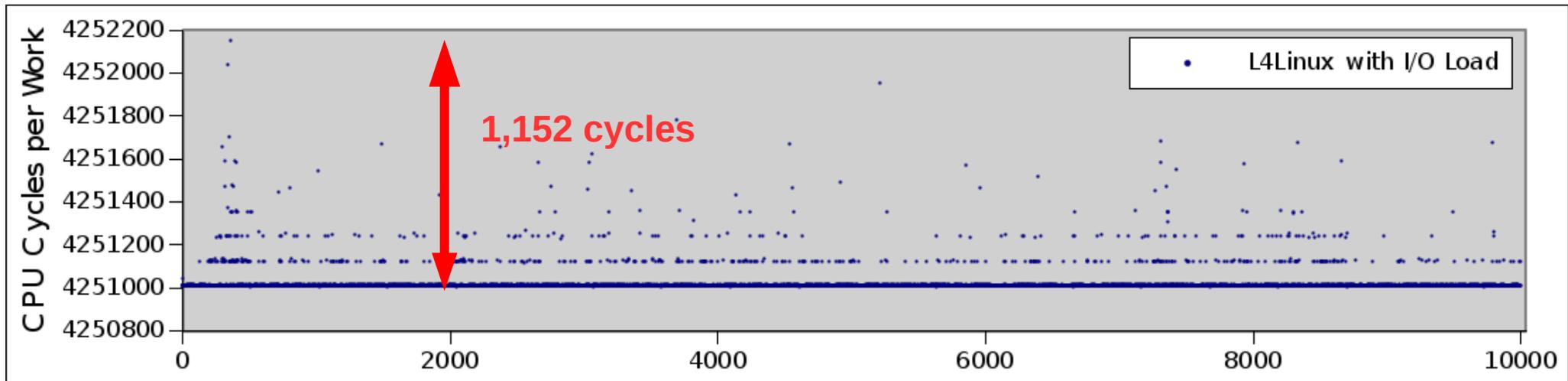
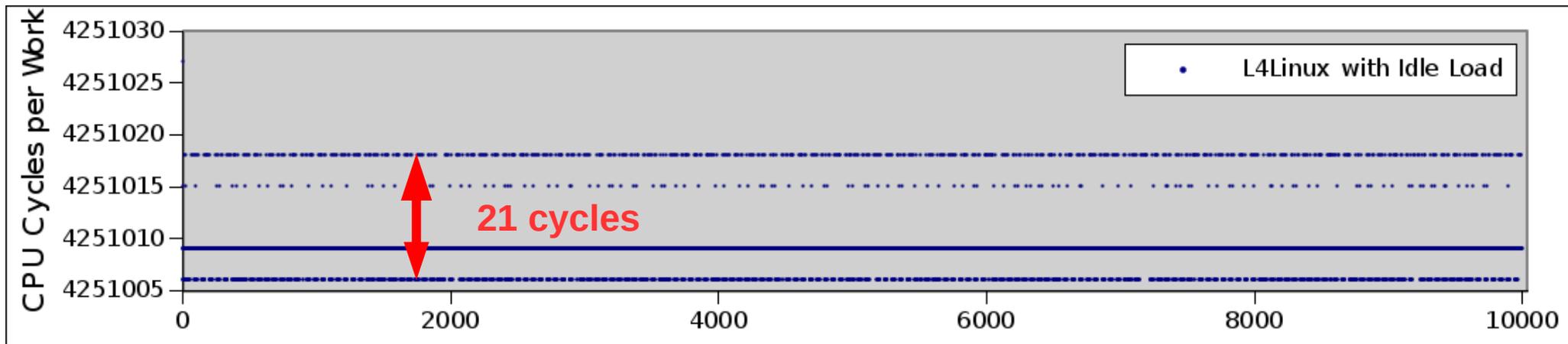
- Linux PREEMPT I/O load



- Detached build load



- With idle, build load & I/O load



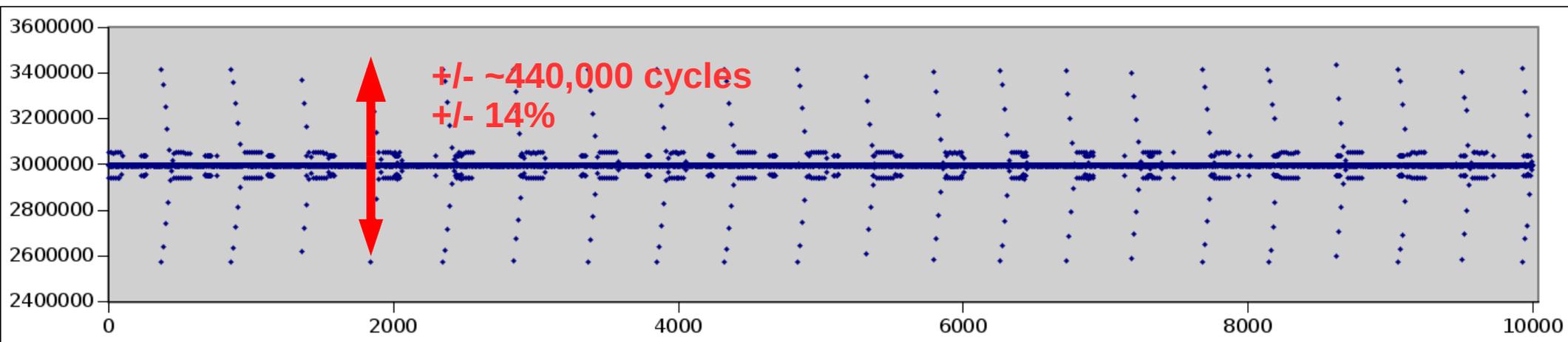
Deltas to minimum in CPU cycles	Linux PREEMPT	Detached Thread in L ⁴ Linux
Idle	434000 (10.2%)	21 (0.0005%)
Build Load	520000 (12.2%)	6528 (0.15%)
I/O Load	773000 (18.2%)	1152 (0.027%)

- Looking at interrupt latency
- Use Linux to launch the following code:

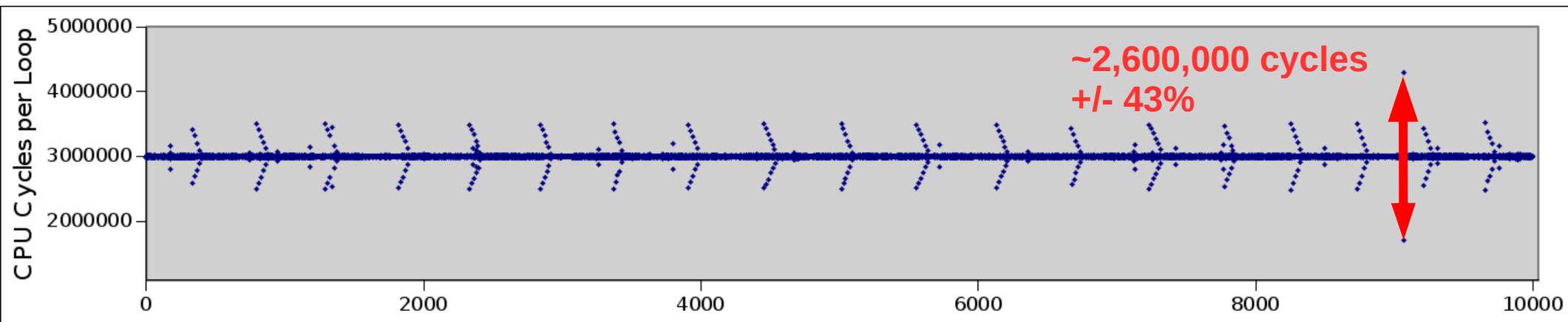
```
next = now() + 1000µs;  
while (1) {  
    wait_for_time(next);  
    tsc = read_tsc();  
    histogramm[i++] = tsc - prev_tsc;  
    prev_tsc = tsc;  
    /* do work */  
    next += 1000µs;  
}
```

- Linux and L4 implementations of `now()` and `wait_for_time()`

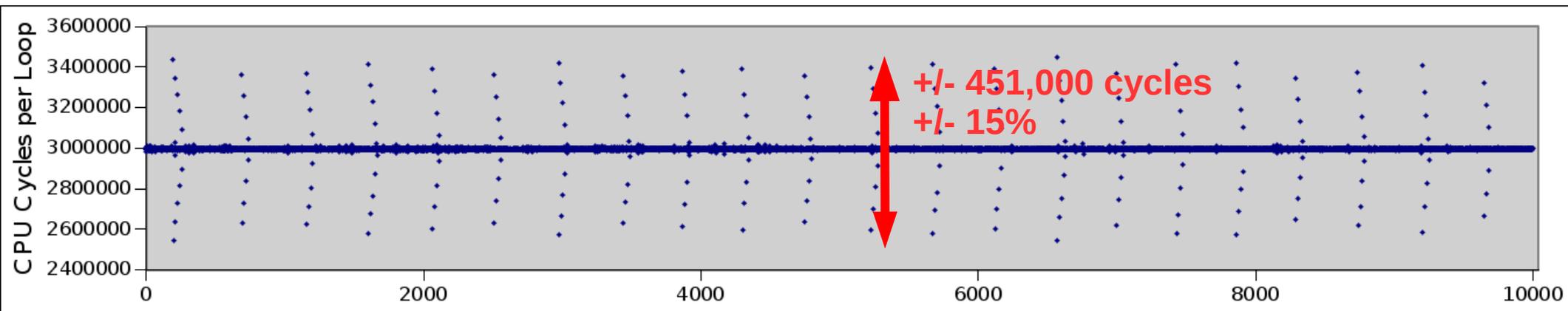
- Linux Idle



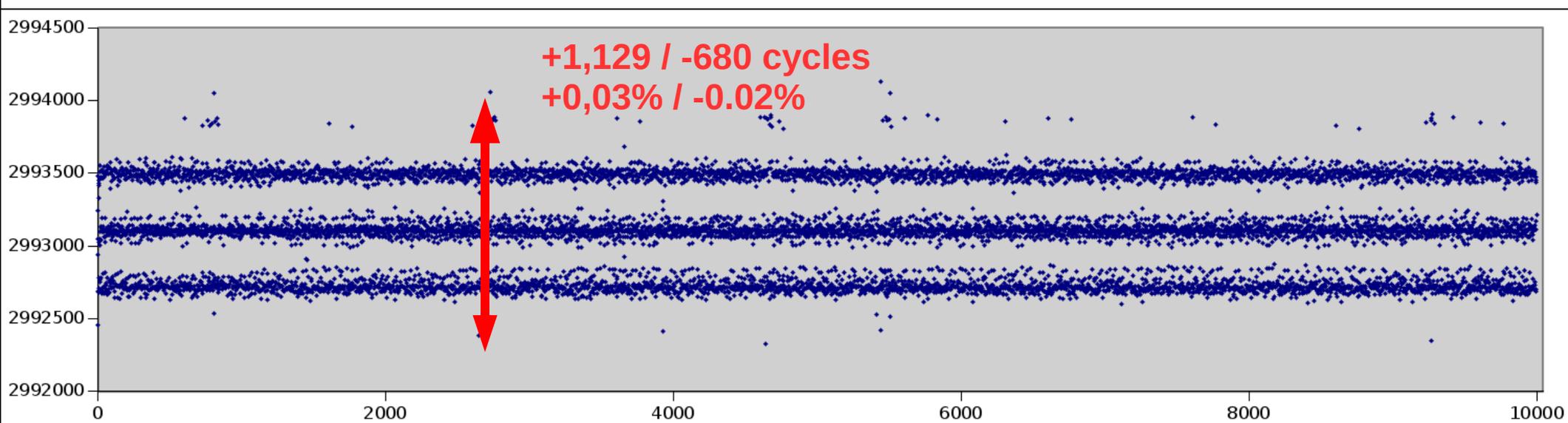
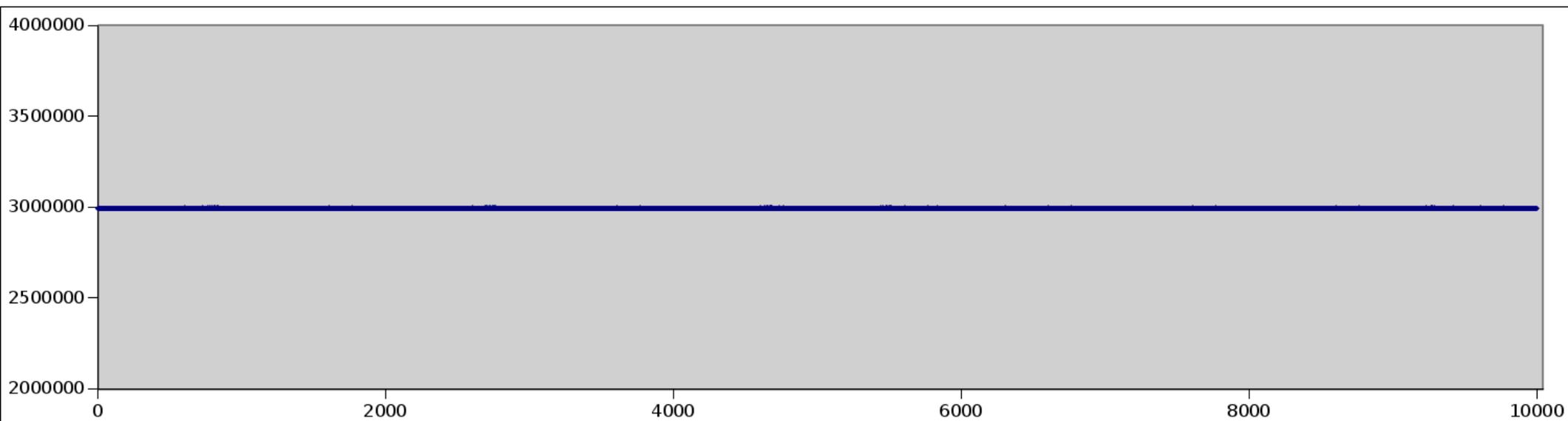
- Linux build load



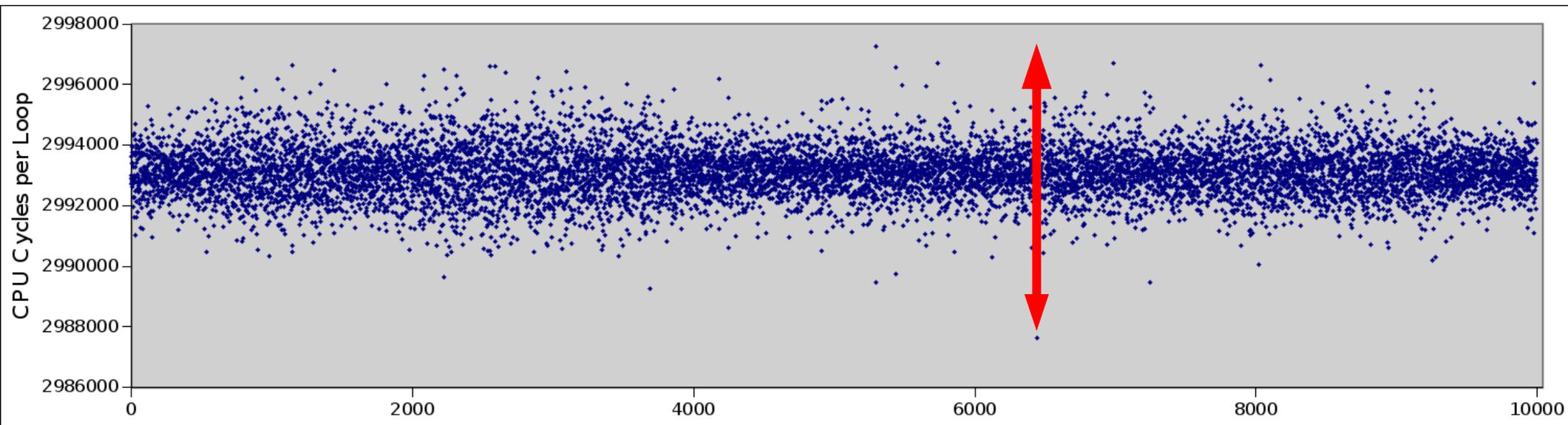
- Linux I/O load



- L⁴Linux detached Idle

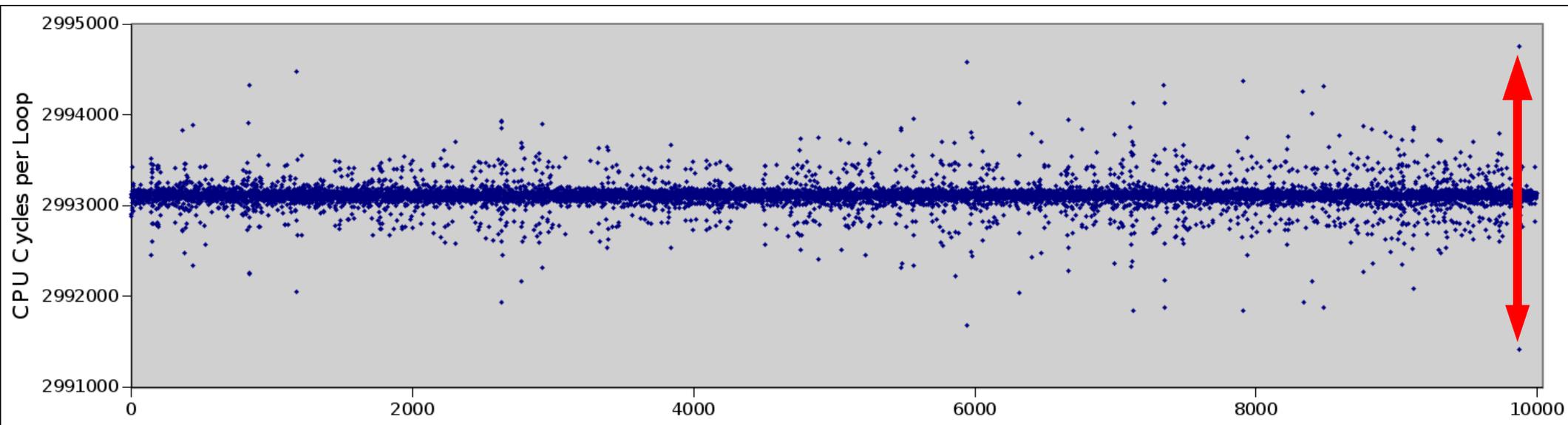


- L⁴Linux detached build load



+4,200 / -5,400 cycles
+0,14% / -0.18%

- L⁴Linux detached I/O load



+1,753 / -1,586 cycles
+0,06% / -0.053%

Maximum Deviation in CPU cycles	Linux	L ⁴ Linux Detached Thread
Idle	862000 (+14,7% / -14.1%)	1809 (+0,03% / -0.02%)
Build Load	2599000 (+/- 43,4)	9633 (+0,14% / -0.18%)
I/O Load	903000 (+/- 15,1%)	3339 (+0,06% / -0.053%)

- Separating a part of a program from a GPOS is possible
- Avoids separation gap
- Real-time and HPC have similar goals

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